

KANGAROOTWELVE: fast hashing based on KECCAK- p

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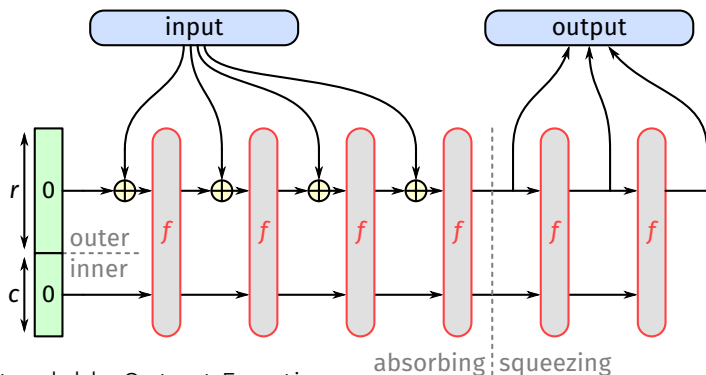
Outline

- 1 What is KANGAROOTWELVE?
- 2 Security vs speed
- 3 Speed vs security

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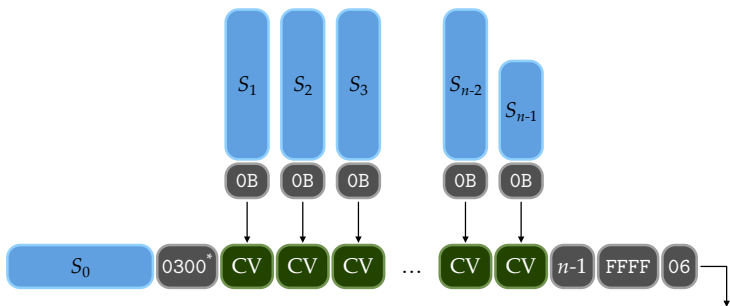
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Let's start from SHAKE128



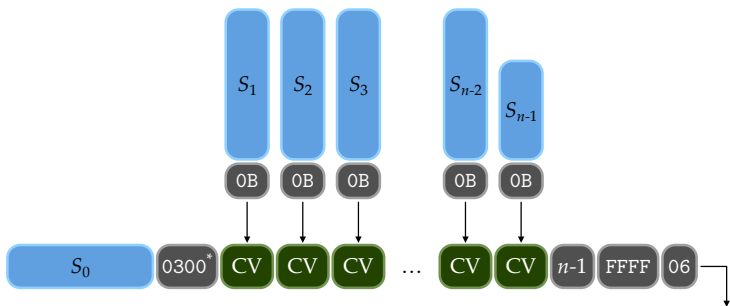
- eXtendable Output Function
- Sponge construction
- Uses KECCAK- $p[1600, n_r = 24]$
- No parallelism at construction level

From SHAKE128 to KANGAROOTWELVE



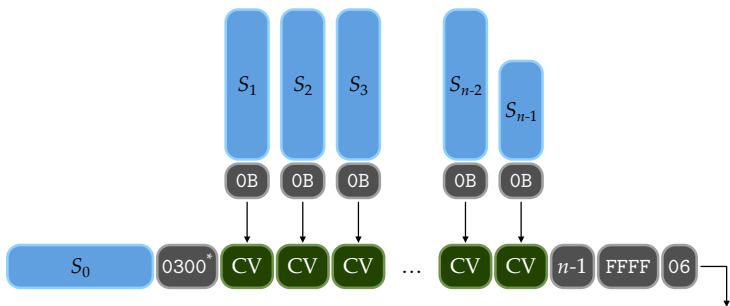
- eXtendable Output Function
- **Tree** on top of sponge construction
- Uses $KECCAK-p[1600, n_r = \mathbf{12}]$
- **Parallelism** grows automatically with input size

From SHAKE128 to KANGAROOTWELVE



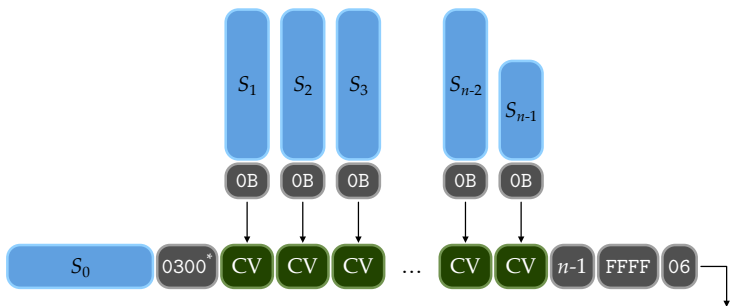
- eXtendable Output Function
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From SHAKE128 to KANGAROOTWELVE



- eXtendable Output Function
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- **Parallelism** grows automatically with input size (per 8KiB)

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KANGAROOTWELVE targets 128-bit security

- Flat sponge claim: 128-bit security strength
 - Collision resistance
 - (Second-) preimage resistance
 - **Multi-target** preimage resistance
 - Chosen-target forced-prefix preimage resistance
 - Correlation-freeness
 - Resistance against length-extension attacks
 - ...
- What about 256-bit security?
 - Philosophically much higher
 - But practically the same: well above the attacker's budget
 - MARSUPILAMIFOURTEEN

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First pillar of security in symmetric cryptography

- Generic security

- Strong mathematical proofs

- ⇒ mode introduces no weaknesses

- ⇒ scope of cryptanalysis focused on primitive

- In our case:

- [EuroCrypt 2008] – On the Indifferentiability of the Sponge Construction

- [IJIS 2014] – Sufficient conditions for sound tree and sequential hashing modes

- [ACNS 2014] – SAKURA: A Flexible Coding for Tree Hashing

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- No proof!

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- Ten years of cryptanalysis on (reduced-round) KECCAK-f[1600]

- ← tune the number of rounds

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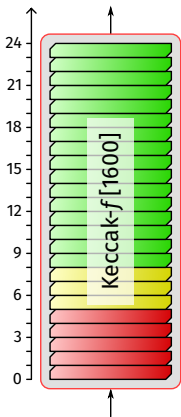
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Status of KECCAK & KANGAROOTWELVE cryptanalysis



- Collision attacks up to 5 rounds
 - Also up to 6 rounds, but for non-standard parameters ($c = 160$)

[Song, Liao, Guo, CRYPTO 2017]
- Distinguishers
 - 7 rounds (practical time)

[Huang et al., EUROCRYPT 2017]
 - 8 rounds (2^{128} time, academic)

[Dinur et al., EUROCRYPT 2015]
- Lots of third-party cryptanalysis available at:

https://keccak.team/third_party.html

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Low-end vs high-end

- How to optimize for both low-end and high-end platforms?
- Avoid 32-bit/64-bit mismatches

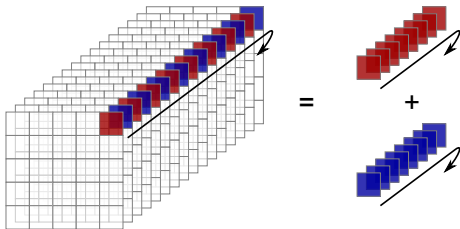
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SHA-256	✓	±
SHA-512	±	✓

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- Avoid 32-bit/64-bit mismatches

	32-bit	64-bit
SHA-256	✓	±
SHA-512	±	✓

Bit interleaving



	32-bit	64-bit
KECCAK- f [800]	✓	±
KECCAK- f [1600]	✓	✓

⇒ let's stick to KECCAK- f [1600]

Exploit parallelism

At the high end:

- SIMD with growing widths
 - 128, 256 and now 512 bits
- Multiple cores

⇒ let's exploit this parallelism

To remain efficient at the low end:

- One-level tree
- Kangaroo hopping

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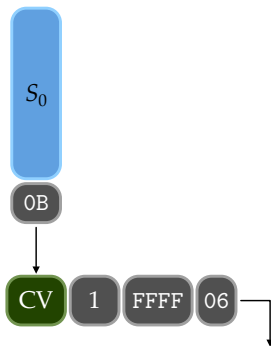
- One-level tree
- Kangaroo hopping

Short messages ($\leq 8\text{KiB}$)

Kangaroo hopping

one call to F

No kangaroo hopping

two calls to F

How fast is KANGAROOTWELVE?

- Twice as fast as SHAKE128 on short inputs $\leq 8\text{KiB}$
- Much faster with parallelism on long inputs $\gg 8\text{KiB}$

	Short input	Long input
Intel® Core™ i5-4570 (Haswell)	3.68 c/b	1.44 c/b
Intel® Core™ i5-6500 (Skylake)	2.89 c/b	1.22 c/b
Intel® Core™ i7-7800X (SkylakeX)	2.06 c/b	0.55 c/b

Single core only.



How fast is KANGAROOTWELVE? (AVX2)

	Skylake (AVX2) cycles/byte
KANGAROOTWELVE	1.22
KANGAROOTWELVE ($\leq 8\text{KiB}$)	2.89
ParallelHash128	2.31
SHAKE128	5.56
SHA-256	6.91
SHA-512	4.64
Blake2bp	1.34
Blake2sp	1.29
Blake2b	3.04
Blake2s	4.85

How fast is KANGAROOTWELVE? (AVX-512)

Skylake (AVX2) vs SkylakeX (AVX-512)

	cycles/byte	
KANGAROOTWELVE	1.22	0.55
KANGAROOTWELVE ($\leq 8\text{KiB}$)	2.89	2.07
ParallelHash128	2.31	0.96
SHAKE128	5.56	4.12
SHA-256	6.91	6.65
SHA-512	4.64	4.44
Blake2bp	1.34	1.39
Blake2sp	1.29	1.22
Blake2b	3.04	2.98
Blake2s	4.85	4.26

Not all optimized for AVX-512 yet

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Any questions?

Thanks for your attention!

- More information

<https://keccak.team/kangarootwelve.html>

- Some implementations

<https://github.com/gvanas/KeccakCodePackage> (C, Python)

<https://github.com/kerukuro/digestpp> (C++)

<https://github.com/mimoo/GoKangarooTwelve> (Go)

<https://rubygems.org/gems/digest-kangarootwelve> (Ruby)

<https://github.com/damaki/libkeccak> (Ada)

- Benoît's RFC draft

<https://datatracker.ietf.org/doc/draft-viguier-kangarootwelve/>